

A MODEL FOR THE INFINITE TENSOR PRODUCT OF GROUPS

INSIGHTS FROM QUANTUM INFORMATION

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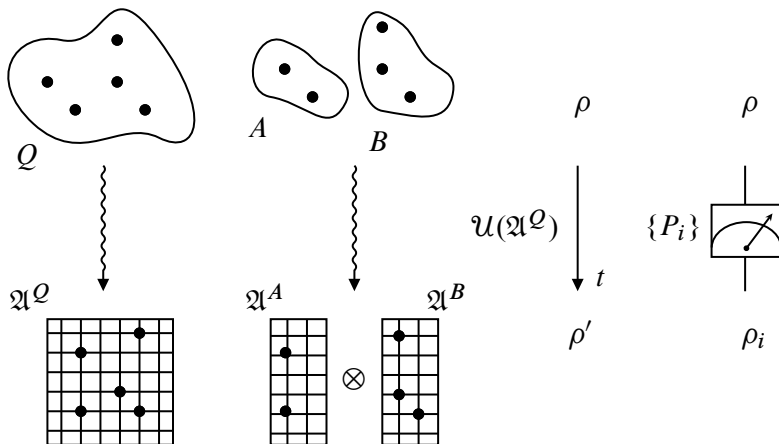
OVERVIEW

- 1 Quantum Setting
- 2 On ∞ -qubit Systems
- 3 Stabilizer Formalism
- 4 Final Remarks

The contemporary mathematical **paradigm** for quantum mechanics can be summarized via four axioms.

The axioms give an account of how a quantum system can be modeled using the rich mathematical framework of **Hilbert spaces**.

QUANTUM AXIOMS VISUALIZED



PAULI GROUP

We call Hilbert spaces $\mathfrak{A} \simeq \mathbb{C}^2$ **qubits**.

Pauli Group

The *Pauli group* \mathcal{P} is the nonabelian matrix group generated by

$$X := \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad Y := \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \quad Z := \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \in \mathbb{M}_2(\mathbb{C}).$$

There is a natural action of \mathcal{P} on a qubit \mathfrak{A} .

If a qubit is “modeled” by the space \mathbb{C}^2 , then n qubits should be modeled by the space

$$\mathbb{C}^2 \underbrace{\otimes \cdots \otimes}_{n \text{ times}} \mathbb{C}^2 = (\mathbb{C}^2)^{\otimes n}.$$

We can generalize the Pauli group to n qubits.

n -QUBIT PAULI GROUP

Denote a **1-local action** of $\Sigma \in \mathcal{P}$ on qubit j of $\mathfrak{A} \simeq \bigotimes_j \mathbb{C}^2$ by

$$\Sigma_j := I_2 \otimes I_2 \otimes \cdots \underbrace{\otimes \Sigma \otimes}_{j \text{th position}} \cdots \otimes I_2.$$

Then, define the n -qubit Pauli group

$$\mathcal{P}_n := \langle iI_j, X_j, Z_j : 1 \leq j \leq n \rangle$$

If $\mathfrak{A} \simeq (\mathbb{C}^2)^{\otimes n}$, then $\Sigma_j \in \mathcal{P}_n$ and $\mathcal{P}_n \curvearrowright \mathfrak{A}$.

IN THE LIMIT

Q: What happens when we take $n \rightarrow \infty$? How could we do algebra in ∞ -qubit systems?

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IN THE LIMIT

$$\left\{ \begin{array}{c} n\text{-qubit systems} \\ \mathcal{P}_n \\ (\mathbb{C}^2)^{\otimes n} \end{array} \right\} \rightsquigarrow \left\{ \begin{array}{c} \infty\text{-qubit systems} \\ \mathcal{P}_\infty \\ \mathfrak{G} \end{array} \right\}$$

FORMAL APPROACH

There are natural inclusions

$$\mathcal{P} = \mathcal{P}_1 \xhookrightarrow{\iota_1^2} \mathcal{P}_2 \xhookrightarrow{\iota_2^2} \dots \xhookrightarrow{\iota_{n-1}^n} \mathcal{P}_n \xhookrightarrow{\iota_n^{n+1}} \dots$$

which send

$$\sigma \mapsto \sigma \otimes I_2.$$

Take the **direct limit** and define $\mathcal{P}_\infty := \operatorname{colim} \langle \mathcal{P}_n, \iota_n^m \rangle$.

CONSTRUCTIVE APPROACH

Consider the invertible 2×2 matrices $\text{GL}_2(\mathbb{C})$. If A_1, \dots, A_n are invertible, then we can form n -fold tensor matrices

$$A_1 \otimes A_2 \otimes \cdots \otimes A_n.$$

We call the group of all these tensors $\text{GL}_2(\mathbb{C})^{\otimes n}$.

CONSTRUCTIVE APPROACH

To generalize, we could take all **sequences** of tensors¹

$$A_{\otimes} := A_1 \otimes A_2 \otimes \cdots \otimes A_i \otimes \cdots$$

and form a group $\overline{\mathrm{GL}_2^{\infty}(\mathbb{C})}$. We call these A_{\otimes} the \mathbb{N} -fold tensor map associated to $\{A_i\}_{i \in \mathbb{N}}$.

¹In practice, it is easier to consider the A_i as automorphisms $\mathbb{C}^2 \xrightarrow{\sim} \mathbb{C}^2$.

CONSTRUCTIVE APPROACH

But doing some algebra tells us we want **finite multilinearity**.
Instead, consider all finitely supported sequences and form the subgroup of restricted N-fold tensor maps $GL_2^\infty(\mathbb{C})$.

Here, all but finitely many of the A_i are trivial.

CONSTRUCTIVE APPROACH

We know \mathcal{P} lives in $\mathrm{GL}_2(\mathbb{C})$, so take all finitely supported tensor sequences from \mathcal{P} and form \mathcal{P}_∞ in $\mathrm{GL}_2^\infty(\mathbb{C})$.

We call this our ∞ -qubit Pauli group.²

²Finding the correct Hilbert space \mathfrak{G} for \mathcal{P}_∞ to act on requires some care. It turns out a natural setting is (isomorphic to) $\ell^2[\Omega]$ for a well-chosen Ω .

ERROR CORRECTION

We model quantum errors as **quantum channels**.

- (i) A *superoperator* is a linear map $\mathcal{E} : \mathbb{B}(\mathfrak{A}) \rightarrow \mathbb{B}(\mathfrak{B})$.
- (ii) A *quantum channel* \mathcal{E} is a superoperator which is completely positive and trace-preserving.

That is, $\mathcal{E} \otimes \text{id}_k \geq 0$ for all k and $\text{tr}(\mathcal{E}\rho) = \text{tr}(\rho)$.

Error Correction

A theorem due to **Karl Kraus** tells us that every quantum channel $\mathcal{E} : \mathbb{B}(\mathfrak{A}) \rightarrow \mathbb{B}(\mathfrak{A})$ has a decomposition³

$$\mathcal{E}(-) = \sum_{x \in X} E_x(-) E_x^\dagger,$$

where $\{E_x\}_{x \in X} \subseteq \mathbb{B}(\mathfrak{A})$.

³Tracking down precisely how large X is and in what topology convergence works can be a bit challenging. If \mathfrak{A} is separable, then $X = \mathbb{N}$ allows ultraweak convergence.

Some terminology:

- (i) A *codespace* is a subspace $\mathcal{C} \subseteq \mathfrak{A}$.
- (ii) Given an error $\mathcal{E} : \mathbb{B}(\mathfrak{A}) \rightarrow \mathbb{B}(\mathfrak{A})$, we call $\mathcal{R} : \mathbb{B}(\mathfrak{A}) \rightarrow \mathbb{B}(\mathfrak{A})$ a *recovery channel* if for all states $\rho \in \mathbb{B}(\mathfrak{A})$,

$$(\mathcal{R} \circ \mathcal{E})(\rho) \propto \rho.$$

on the codespace.

- (iii) An error \mathcal{E} is *correctable* if a codespace \mathcal{C} and recovery channel \mathcal{R} exist.

GOTTESMAN'S STABILIZER FORMALISM

Theorem (Gottesman, 1997)

Let $\mathfrak{A} \simeq (\mathbb{C})^{\otimes n}$, let \mathcal{E} be a noise channel on \mathfrak{A} with Kraus operators $\{E_i\}_{i=1}^r$, and let $\mathcal{S} \leq \mathcal{P}_n$ be a finitely-generated abelian subgroup without $-(I^{\otimes n})$. Then, \mathcal{E} is correctable on $\text{Fix}_{\mathcal{S}}(\mathfrak{A})$ if and only if for all $1 \leq i, j \leq r$, we have

$$E_i^\dagger E_j \in \text{span}\{(\mathcal{P}_n \setminus \mathcal{N}_{\mathcal{P}_n}(\mathcal{S})) \cup \mathcal{S}\} \subseteq \mathbb{B}(\mathfrak{A}),$$

where $\mathcal{N}_{\mathcal{P}_n}(\mathcal{S})$ is the normalizer of \mathcal{S} in the n -qubit Pauli group.

∞ -QUBIT STABILIZER FORMALISM

Letting \mathfrak{G} be an ∞ -qubit space, can we recover the Stabilizer Formalism for errors on \mathfrak{G} by simply “exchanging n with ∞ ?”

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