

# COMMUTANT OF THE CLIFFORD GROUP

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REPRESENTATION-THEORETIC METHODS IN QIT

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# OVERVIEW

## 1 Preliminaries and Motivation

## 2 Computing the Commutant

## 3 Applications

## 4 References

**Q:** Why should we be interested in computing the commutant of the Clifford group?

**A:** The image of Clifford twirling is contained in the commutant.

## PAULI GROUP

The Pauli matrices are

$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \quad \text{and} \quad Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}.$$

In  $\mathcal{U}_2$ , they generate the 1-qubit Pauli group  $\mathcal{P}_1$ . Taking all  $n$ -fold tensor products

$$\sigma_1 \otimes \sigma_2 \otimes \cdots \otimes \sigma_n, \quad \sigma_j \in \mathcal{P}_1,$$

we obtain the  $n$ -qubit Pauli group.



## PAULI GROUP

Quotienting by phases  $\langle iI \rangle$  yields a group

$$\mathcal{P}_n = \{I, X, Y, Z\}^{\otimes n}.$$

Denote all  $m$ -tuples of operators in  $\mathcal{P}_n$  by  $\mathcal{P}_n^m$ .

# STABILIZER CODES

The fixed subspace of abelian subgroups (without  $-I^{\otimes n}$ ) of the  $n$ -qubit Pauli group give rise to stabilizer codes—the quintessential sort of error-correcting code.

# CLIFFORD GROUP

The Clifford group  $\mathcal{C}_n$  is the subgroup of  $\mathcal{U}_{2^n}$  given by

$$\mathcal{C}_n = \left\{ C \in \mathcal{U}_{2^n} : C\sigma C^\dagger \in \mathcal{P}_n \text{ for all } \sigma \in \mathcal{P}_n, \text{ up to phases} \right\}.$$

# RANDOM CLIFFORD GATES

To study the effects of applying random Clifford gates to a system, we characterize averaging over the Clifford group using  $k$ -fold twirling.

The Haar measure allows us to define  $k$ -fold twirling for any subgroup  $G \subseteq \mathcal{U}_{2^n}$ . We focus on  $\mathcal{C}_n$ .

Let  $k$  be a number,  $G \subseteq \mathcal{U}_{2^n}$  be a subgroup, and  $\mathcal{K}$  be  $k$  copies of an  $n$ -qubit Hilbert space (i.e., of dimension  $2^{nk}$ ).

## COMMUTANT

The  $k$ -fold commutant of  $G$  is the  $\mathbb{C}$ -algebra

$$\text{Com}(G^{\otimes k}) = \left\{ \rho \in \mathbb{B}(\mathcal{K}) : U^{\dagger \otimes k} \rho U^{\otimes k} = \rho \text{ for all } U \in G \right\}.$$

# CLIFFORD TWIRLING

The Clifford twirling (super)operator is

$$\Phi_{\text{Cl}}^{(k)}(-) = \frac{1}{|\mathcal{C}_n|} \sum_{C \in \mathcal{C}_n} C^{\otimes k}(-) C^{\dagger \otimes k}.$$

# CLIFFORD TWIRLING

## Lemma

Let  $U \in \mathcal{C}_n$  be an arbitrary Clifford operator. Then, for all  $X \in \mathbb{B}(\mathcal{K})$

$$U^{\dagger \otimes k} \Phi_{\text{Cl}}^{(k)}(X) U^{\otimes k} = \Phi_{\text{Cl}}^{(k)}(X).$$

That is, the image  $\Phi_{\text{Cl}}^{(k)}(\mathbb{B}(\mathcal{K})) \subseteq \text{Com}(\mathcal{C}_n^{\otimes k})$ .

## CLIFFORD TWIRLING

## Proof.

For any Clifford operator  $C \in \mathcal{C}_n$ ,

$$U^{\dagger \otimes k} C^{\otimes k} X C^{\dagger \otimes k} U^{\otimes k} = (U^{\dagger} C)^{\otimes k} X (U^{\dagger} C)^{\dagger \otimes k}.$$

Left multiplication  $U^{\dagger} \times - : \mathcal{C}_n \rightarrow \mathcal{C}_n$  is a bijection, so we have just reindexed the twirling sum. □

## TWIRLING

In fact, the same holds for twirling (via the Haar measure) over any group  $G \subseteq \mathcal{U}_{2^n}$ . Just use the invariance properties.

## BASIS

If we can find a nice basis for  $\text{Com}(\mathcal{C}_n^{\otimes k})$ , we can then write down any averaged Clifford operator in terms of this basis.

## BASIS

In [GNW21], a basis is found when  $k \leq n + 1$ . This basis is built from operators of the form

$$r(T) = \sum_{(x,y) \in T} |x\rangle\langle y|,$$

where  $T \subseteq \mathbb{F}_2^k \oplus \mathbb{F}_2^k$  is a *stochastic Lagrangian subspace* with respect to a symplectic form (see later on Robust Hudson Theorem).



## DIMENSION

This construction tells us

$$\dim \left( \text{Com}(\mathcal{C}_n^{\otimes k}) \right) = \prod_{j=0}^{k-2} (2^j + 1).$$

That is, when  $k \leq n + 1$ , the dimension is independent of  $n$ .

## ADJACENCY MATRIX

Let  $\Gamma = (V, E)$  be a finite graph with  $m$  vertices  $v_1, \dots, v_m$ .

The adjacency matrix  $\mathcal{A}(\Gamma)$  is an operator in  $\mathbb{M}_m(\mathbb{F}_2)$  defined by

$$\mathcal{A}(\Gamma)_{i,j} = \begin{cases} 1, & (v_i, v_j) \in E \\ 0, & \text{otherwise.} \end{cases}$$

## ANTICOMMUTATION GRAPH

Let  $\mathbf{P} = (P_1, P_2, \dots, P_m) \in \mathcal{P}_n^m$ . The anticommutation graph of  $\mathbf{P}$  has

- (i)  $m$  vertices  $P_1, \dots, P_m$ .
- (ii) edges between anticommuting pairs  $(P_i, P_j)$ .

Write  $\mathcal{A}(\mathbf{P})$  for the adjacency matrix of the above graph.



Now, a useful fact:

Lemma ([BEL<sup>+</sup>25], Lemma 17)

Let  $\mathbf{P} \in \mathcal{P}_n^m$  be algebraically independent. Let  $\mathbf{Q} \in \pm \mathcal{P}_n^m$  be algebraically independent with signs. Then, there exists a Clifford  $C \in \mathcal{C}_n$  such that

$$\mathbf{C}\mathbf{P}\mathbf{C}^\dagger = \mathbf{Q}.$$

if and only if

$$\mathcal{A}(\mathbf{P}) = \mathcal{A}(\mathbf{Q}).$$



## Idea.

When  $m = 2n$ , the Clifford operator is given by

$$C(X) = \frac{1}{2} \sum_{\alpha \in \mathbb{F}_2^m} \text{tr} \left( P_\alpha^\dagger X \right) Q_\alpha,$$

where

$$P_\alpha = \prod_{j=1}^m P_j^{\alpha_j}, \quad \alpha \in \mathbb{F}_2^m$$

and likewise for  $Q_\alpha$ . □

## Idea, continued.

If  $m < 2n$ , we can complete both  $\mathbf{P}$  and  $\mathbf{Q}$  with  $2n - m$  algebraically independent Pauli operators while keeping the adjacency matrices agreeing.

That we can do this relies on the canonical graph form of any anticommutation graph. □

## INDEPENDENT GRAPH-BASED PAULI MONOMIALS

## Definition

Let  $G \in \text{Sym}_0(\mathbb{F}_2^{m \times m})$  be a symmetric binary matrix, and  $V \in \text{Even}(\mathbb{F}_2^{k \times m})$ . We define independent graph-based Pauli monomials as

$$\mathcal{U}_I([V, G]) = \frac{1}{|S_{[V, G]}|} \sum_{\substack{\mathbf{P} \in \mathcal{P}_n^{\times m} \\ \mathcal{A}(\mathbf{P}) = G, \\ \mathbf{P} \text{ alg ind.}}} \prod_{j=1}^m P_j^{\otimes v_j}$$

i.e., the sum runs over the set of algebraically independent Pauli operators  $\mathbf{P} = (P_1, \dots, P_m)$ .

# ORTHOGONAL BASIS OF THE COMMUTANT

## Theorem

*The set of independent graph-based Pauli monomials  $\{\mathcal{U}_I([V, G])\}_{V \in \text{Even}(\mathbb{F}_2^{k \times m}), G \in \text{Sym}_0(\mathbb{F}_2^{m \times m})}$  forms an orthogonal basis of the commutant  $\text{Com}(\mathcal{C}_n^{\otimes k})$ .*

*Proof Sketch.* The commutant is precisely the image of the twirling map

$$\Phi_{\text{Cl}}^{(k)}(P) = \frac{1}{|\mathcal{C}_n|} \sum_{C \in \mathcal{C}_n} C^{\otimes k} P C^{\dagger \otimes k}.$$

Now it can be shown that  $V$  and  $G$  are invariant under the action of  $C$ , so we can combine the monomials with equivalent  $V$  and  $G$  together.



# PAULI MONOMIALS

## Definition

Let  $M \in \text{Sym}_0(\mathbb{F}_2^{m \times m})$ , and  $V \in \text{Even}(\mathbb{F}_2^{k \times m})$ . The Pauli monomial, denoted by  $\Omega(V, M)$  is defined as

$$\Omega(V, M) := \frac{1}{d^m} \sum_{P \in \mathcal{P}_n^m} P_1^{\otimes v_1} P_2^{\otimes v_2} \cdots P_m^{\otimes v_m} \times \left( \prod_{\substack{i, j \in [m] \\ i < j}} \chi(P_i, P_j)^{M_{i,j}} \right),$$

## GRAPHICAL CALCULUS

$$\Omega \left( \begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 1 & 1 \end{pmatrix}, \begin{pmatrix} 0 & 0 & 1 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix} \right) = \begin{array}{c} \text{A 6x3 grid of dots with a wavy line underneath. The first column has 6 dots, the second has 3, and the third has 1.} \\ \text{The first column is highlighted with a wavy line.} \end{array}$$

Figure: Graphical representation

$$\begin{array}{c} \text{A 6x3 grid of dots with a wavy line underneath. The first column has 6 dots, the second has 3, and the third has 1.} \\ \text{The first column is highlighted with a wavy line.} \end{array} = \begin{array}{c} \text{A 6x3 grid of dots with a wavy line underneath. The first column has 6 dots, the second has 3, and the third has 1.} \\ \text{The first column is highlighted with a wavy line.} \end{array}$$

Figure: Adding first column to the second

# GRAPH-BASED PAULI MONOMIALS

## Definition

Let  $G \in \text{Sym}_0(\mathbb{F}_2^{m \times m})$ , and  $V \in \text{Even}(\mathbb{F}_2^{k \times m})$ . The graph-based Pauli monomials is defined as

$$U(V, G) := \frac{1}{d^m} \sum_{\substack{P \in \mathcal{P}_n \\ \mathcal{A}(P)=G}} \prod_{j=1}^m P_j^{\otimes v_j},$$

where  $\mathcal{A}(P)$  denotes the adjacency matrix of the anticommutation graph associated with the set  $\{P_1, \dots, P_m\}$ .



# Connection Between Different Bases

Pauli monomials  $\Omega(V, M)$  and graph-based Pauli monomials  $\mathcal{U}(V, G)$  are related via a linear Fourier transform:

$$\Omega(V, M) = \sum_{G \in \text{Sym}_0(\mathbb{F}_2^{m \times m})} (-1)^{\sum_{i < j} M_{i,j} G_{i,j}} \mathcal{U}(V, G).$$

Also there is a bijection between graph-based Pauli monomials and independent graph-based Pauli monomials.

## Corollary

*The sets  $\{\Omega(V, M)\}_{V, M}$ ,  $\{\mathcal{U}(V, G)\}_{V, G}$ , and  $\{\mathcal{U}_I(V, G)\}_{V, G}$  generate the same vector space.*

# CLIFFORD-WEINGARTEN CALCULUS

## Theorem

Let  $\Phi_{\text{Cl}}^{(k)}(O)$  be the  $k$ -fold twirling operator  $O \in \mathcal{B}(\mathcal{K})$  and let  $\mathcal{P}$  be the set of (reduced) Pauli monomials. Then it reads

$$\Phi_{\text{Cl}}^{(k)}(O) = \sum_{\Omega, \Omega' \in \mathcal{P}} (\mathcal{W}^{-1})_{\Omega, \Omega'} \text{tr}(O\Omega)\Omega'$$

where we call the coefficients  $(\mathcal{W}^{-1})_{\Omega, \Omega'}$  Clifford-Weingarten functions, which can be obtained as the (pseudo-)inverse of the Gram-Schmidt matrix  $\mathcal{W}_{\Omega, \Omega'} := \text{tr}(\Omega^\dagger \Omega')$ .

# CLIFFORD-WEINGARTEN CALCULUS

## Theorem

Let  $\mathcal{W}$  the Gram-Schmidt matrix, and let  $\mathcal{W}^{-1}$  be its inverse. Let  $n \geq k^2 - 3k + 7$ . Then the properties

$$\left| (\mathcal{W}^{-1})_{\Omega, \Omega} - \frac{1}{d^k} \right| \leq \frac{6|\mathcal{P}|^2}{d^{k+1}},$$

$$|(\mathcal{W}^{-1})_{\Omega, \Omega'}| \leq \frac{5|\mathcal{P}|^2}{d^{k+1}},$$

hold true.

## DE FINETTI THEOREM FOR STABILIZER SYMMETRIES

Recall the regular de Finetti theorem:

### Theorem

Let  $\rho$  be a quantum state on  $(\mathbb{C}^\ell)^{\otimes k}$  that commutes with all permutations. Then, there exists a probability measure  $\mu$  on the space of mixed states on  $\mathbb{C}^\ell$  such that

$$\frac{1}{2} \left\| \rho_{1\dots s} - \int d\mu(\sigma) \sigma^{\otimes s} \right\|_1 \leq 2\ell^2 \frac{s}{k}$$

## DE FINETTI THEOREM FOR STABILIZER SYMMETRIES

Let's add symmetries. For qudits with  $d$  being prime, consider the additional group of symmetries  $O_k(d)$  consisting of  $k \times k$ -matrices with entries in  $\mathbb{Z}_d$  which are orthogonal and such that the sum of elements in each row is equal to 1 (mod  $d$ ). We call this group *stochastic orthogonal group*.

## Theorem

Let  $d$  be a prime and  $\rho$  a quantum state on  $((\mathbb{C}^d)^{\otimes n})^{\otimes k}$  that commutes with the action of  $O_k(d)$ . Then, there exists a probability distribution  $p$  on the (finite) set of mixed stabilizer states of  $n$  qudits, such that

$$\frac{1}{2} \left\| \rho_{1\dots s} - \sum_{\sigma_S} p(\sigma_S) \sigma_S^{\otimes s} \right\|_1 \leq 2d^{\frac{1}{2}(2n+2)^2} d^{-\frac{1}{2}(k-s)}.$$

## ROBUST HUDSON THEOREM

Let  $x = (p, q) \in \mathbb{Z}^{2n}$ . Define the *Weyl operator*

$$W_x = -\tau^{-p \cdot q} (Z^{p_1} X^{q_1}) \otimes \cdots \otimes (Z^{p_n} X^{q_n}).$$

Consider a symplectic form on  $\mathbb{Z}^{2n}$  by  $[x, x'] = p \cdot q' - q \cdot p'$ . Then for any operator  $B$  on  $(\mathbb{C}^d)^{\otimes n}$  one can define the *Wigner function*  $w_B: \mathbb{Z}_d^{2n} \rightarrow \mathbb{C}$  by

$$w_B(x) = d^{-2n} \sum_y \omega^{-[x,y]} \text{tr}(W_y^\dagger B)$$

where  $\omega = e^{2\pi i/d}$ .

# ROBUST HUDSON THEOREM

For odd  $d$ , the Wigner function of a quantum state is real and  $-d^{-n} \leq w_\psi(x) \leq d^{-n}$ . Furthermore the Wigner function is a quasi-probability distribution, i.e.,  $\sum_x w_\psi(x) = 1$ .

# ROBUST HUDSON THEOREM

## Theorem

*A pure state  $|\psi\rangle$  is a stabilizer state if and only if it has a nonnegative Wigner function.*

For a quantitative version, define *sum-negativity* of  $\psi$  to be

$$\text{sn}(\psi) = \sum_{w_\psi(x) < 0} |w_\psi(x)|.$$



# ROBUST HUDSON THEOREM

## Theorem

*Let  $d$  be odd and  $\psi$  a pure quantum state of  $n$  qudits. Then there exists a stabilizer state  $|S\rangle$ , such that  $|\langle S|\psi\rangle|^2 \geq 1 - 9d^2\text{sn}(\psi)$*

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