

STABILIZER FORMALISM IN ∞ -QUBIT SYSTEMS

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OVERVIEW

- 1 Errors as Channels
- 2 Stabilizer Formalism
- 3 Extension to Spin Chains
- 4 Locality Two Ways
- 5 Final Remarks

Quantum computation is inherently susceptible to **quantum noise**—families of potential errors triggered by quantum mechanical phenomena.

So, what is noise?

Knill and Laflamme say: a (quantum) channel!

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Knill and Laflamme say: a (quantum) **channel!**

Alright, then what is a quantum channel?

Fix a state space \mathfrak{A} .

Maybe a qubit \mathbb{C}^2 .

Or better yet, maybe n qubits

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A **channel** on \mathfrak{A} is a bounded linear map on $\mathcal{S}_1(\mathfrak{A})$ of a specific form.

Start with the bounded operators $\mathfrak{A} \rightarrow \mathfrak{A}$ of finite trace.¹

$$\mathfrak{S}_1(\mathfrak{A})$$

¹This is a Banach space.

Couple with the environment \mathfrak{B} .

$$\begin{array}{c} \mathfrak{S}_1(\mathfrak{A}) \\ \downarrow -\otimes\sigma \\ \mathfrak{S}_1(\mathfrak{A} \otimes \mathfrak{B}) \end{array}$$

Evolve unitarily.

$$\begin{array}{c} \mathfrak{S}_1(\mathfrak{A}) \\ \downarrow -\otimes\sigma \\ \mathfrak{S}_1(\mathfrak{A} \otimes \mathfrak{B}) \\ \downarrow U(-)U^\dagger \\ \mathfrak{S}_1(\mathfrak{A} \otimes \mathfrak{B}) \end{array}$$

Trace off the environment.

$$\begin{array}{c} \mathfrak{S}_1(\mathfrak{A}) \\ \downarrow - \otimes \sigma \\ \mathfrak{S}_1(\mathfrak{A} \otimes \mathfrak{B}) \\ \downarrow U(-)U^\dagger \\ \mathfrak{S}_1(\mathfrak{A} \otimes \mathfrak{B}) \\ \downarrow \text{tr}_{\mathfrak{B}} \\ \mathfrak{S}_1(\mathfrak{A}) \end{array}$$

Every quantum channel \mathcal{E} on \mathfrak{A} decomposes as a sum¹

$$\mathcal{E}(\rho) = \sum_k E_k \rho E_k^\dagger, \quad E_k \in \mathbb{B}(\mathfrak{A}).$$

We call the E_k the **Kraus** (or error) operators of \mathcal{E} .

¹If $\dim(\mathfrak{A})$ is infinite, this series converges in trace norm.

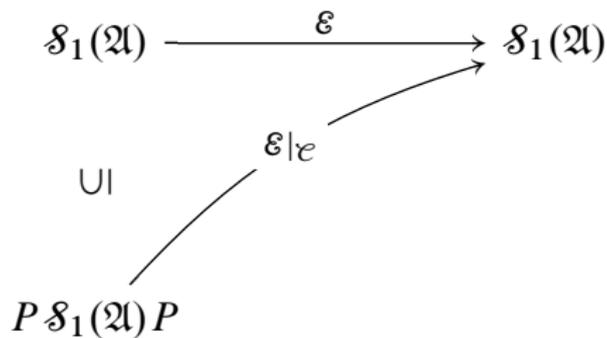
Say \mathcal{C} is a (closed) subspace $P : \mathfrak{A} \rightarrow \mathcal{C} \subseteq \mathfrak{A}$ with orthonormal basis $(|j\rangle)_j$. The **Knill-Laflamme** correctability condition tells us that a channel \mathcal{E} has a retraction channel \mathcal{R} over \mathcal{C} if and only if

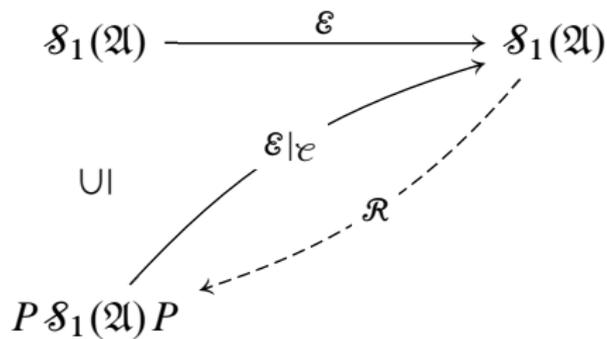
$$\langle j | E_\ell^\dagger E_m | j \rangle = \langle k | E_\ell^\dagger E_m | k \rangle$$

and

$$\langle j | E_\ell^\dagger E_m | k \rangle = 0,$$

for all j, k, ℓ, m .





The **Pauli group** \mathcal{P} on 1 qubit is the order-16 matrix group generated by the Pauli spin matrices.

$$\mathcal{P} = \langle i \rangle \{ \mathbf{1}, X, Y, Z \} \subseteq \text{GL}_2(\mathbb{C})$$

The n -qubit Pauli group \mathcal{P}_n consists of all n -fold (Kronecker) tensor products of elements of \mathcal{P} .

$$\mathcal{P}_n = \{\sigma_1 \otimes \cdots \otimes \sigma_n : \sigma_j \in \mathcal{P}\} \subseteq \text{GL}_{2^n}(\mathbb{C})$$

Gottesman's stabilizer formalism lets us do error correction algebraically using the Pauli group.

If \mathcal{E} is an n -qubit channel whose error operators are all from \mathcal{P}_n , then \mathcal{E} is correctable on the \mathcal{S} -fixed² subspace $\text{Fix}_{\mathcal{S}}(\mathcal{A})$ precisely when

$$E_k^\dagger E_\ell \in (\mathcal{P}_n \setminus C_{\mathcal{P}_n}(\mathcal{S})) \cup \mathcal{S}.$$

²Here, \mathcal{S} is an abelian subgroup without the negative identity $-\mathbb{1}_2 \otimes \cdots \otimes \mathbb{1}_2$.

Suppose our physical system is a countable collection of coupled qubits.
For instance, a chain³ of spin-1/2 particles.

³For convenience, we index the sites by $\mathbb{N} = 1, 2, 3, \dots$
instead of the more usual choice \mathbb{Z} .

Then, our state space should be an infinite tensor product

$$\mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \dots$$

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Instead, we say an ∞ -qubit Hilbert space \mathfrak{H} is spanned by infinite sequences of the (standard) basis of \mathbb{C}^2 , where we allow finite variation from a chosen basis vector in each copy of \mathbb{C}^2 .

Define the ∞ -qubit Pauli group \mathcal{P}_∞ be the set of tensor sequences of operators in \mathcal{P} , where all but finitely many operators are trivial.

$$\mathcal{P}_\infty = \{\sigma_1 \otimes \sigma_2 \otimes \cdots : \sigma_j \in \mathcal{P} \text{ and } |\{\sigma_j \neq \mathbb{1}_2\}| < \infty\} \subseteq \mathbb{B}(\mathfrak{H}).$$

Replacing all the n s in the stabilizer formalism with ∞ s, we obtain a stabilizer formalism for ∞ -qubit systems!

We now ask, “what does \mathcal{P}_∞ look like when we zoom in?”

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Observe that there is a sequence of inclusions

$$\mathcal{P} \hookrightarrow \mathcal{P}_2 \hookrightarrow \mathcal{P}_3 \hookrightarrow \dots \hookrightarrow \mathcal{P}_n \hookrightarrow \dots$$

where we send

$$\sigma \mapsto \sigma \otimes \mathbb{1}_2$$

in every map.

We are working with groups, so we can take the **direct limit**

$$\operatorname{colim}_n(\mathcal{P} \hookrightarrow \mathcal{P}_2 \hookrightarrow \mathcal{P}_3 \hookrightarrow \dots \hookrightarrow \mathcal{P}_n \hookrightarrow \dots).$$

This direct limit is exactly (isomorphic to) \mathcal{P}_∞ !

Colimits like this come equipped with embeddings

$$\xi_n : \mathcal{P}_n \hookrightarrow \mathcal{P}_\infty,$$

given here by placing the n -qubit system in the n th open slot.

If $n \neq m$, then

$$[\xi_n(\Sigma), \xi_m(\Sigma)] = 0,$$

so \mathcal{P}_∞ exhibits locality.

But, in what sense is \mathcal{P}_∞ a tensor product of infinitely many copies of \mathcal{P} ?

For that matter, in what sense is \mathcal{P}_n a finite tensor power of \mathcal{P} ?

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This language lets us define a tensor product operation

$$\otimes : \text{Grp}(C_4) \times \text{Grp}(C_4) \rightarrow \text{Grp}(C_4),$$

constructed using a good notion of **C_4 -bilinearity**.

Using this C_4 -tensor product, we can say that

$$\mathcal{P}_n \simeq \mathcal{P} \underbrace{\otimes \cdots \otimes}_{n} \mathcal{P}$$

as C_4 -groups.

Then, we can extend this to an isomorphism

$$\mathcal{P}_\infty \simeq \mathcal{P}^{\otimes \infty},$$

where $\mathcal{P}^{\otimes \infty}$ is the (finitely supported) infinite C_4 -tensor power of \mathcal{P} .

So not only does the ∞ -qubit Pauli group behave n -locally like the finite Pauli groups, but **1-locally** too!

Most things we have done here are rather general.

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Further reading:

- (L) Felix Leditzky's notes *Quantum Channels*.
- (NC) Nielsen and Chuang's book *Quantum Computation and Quantum Information*, sections §8 and §10.
- (KLM) Kaye, Laflamme, and Mosca's book *An Introduction to Quantum Computing*, section §10.

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Thanks! Questions?

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